# **GRAPH-THEORETICAL OPTIMIZATION OF FUSION-BASED GRAPH STATE GENERATION**

# **Seok-Hyung Lee1,2 and Hyunseok Jeong1**

1Department of Physics and Astronomy, Seoul National University, Seoul, Republic of Korea 2Centre for Engineered Quantum Systems, School of Physics, University of Sydney, Sydney, NSW, Australia

### **SUMMARY**

- ‣ **Graph states**: Resources for various quantum information task including measurement-based quantum computing, fusion-based quantum computing, quantum repeaters, and quantum metrology
- ‣ In linear optical systems, graph states can be generated by **type-II fusion operations** → **Non-deterministic**, thus large graph states are difficult to generate.
- ‣ We propose **a graph-theoretical strategy to effectively optimize the generation of any graph state via type-II fusions.**
- ‣ Main idea: find a **graph state equivalent to the desired graph state** under local Clifford gates and type-II fusions but easier to generate → **Unraveling**

 $\rightarrow$  Used in measurement-based quantum computing [1,2], fusion-based quantum computing [3], quantum repeaters [4], quantum metrology [5], etc.

‣ We expect that our strategy and software will assist researchers in developing and accessing experimentally viable schemes that use photonic graph states.

# **BACKGROUND**

 $\blacktriangleright$  **Graph state:** For a graph  $G = (V, E)$ ,  $|G\rangle_V := \prod U_{CZ} (v_1, v_2) \bigotimes |+\rangle_V$ *v*<sub>1</sub>,*v*<sub>2</sub>∈*E v*∈*V*  $\forall v \in V$ ,  $S_v | G \rangle := | X_v |$ *u*∈adj(*v*)  $Z_u \mid |G\rangle = |G\rangle$ 



‣ Equivalence of graph states under local Clifford gates:

 $\exp\left[-i\frac{1}{4}X_{\nu}\right] - \prod_{i=1}^{\infty}\exp\left[i\frac{1}{4}Z_{\nu}\right] \mid G\rangle = |{\rm LC}_{\nu}(G)\rangle \to {\sf Local\ complementation}$  (LC) *π* 4  $X_{\scriptscriptstyle\mathcal{V}}^{}$  $\|$   $\|$ *u*∈adj(*v*)  $\exp |i \rangle$ *π* 4  $Z_u | G \rangle = | LC_v(G) \rangle$ 

→ Connect/disconnect every pair of adjacent vertices of two vertices

### **OVERVIEW OF THE STRATEGY**

- ‣ Basic resource state: **three-qubit linear graph state** |*G*(3) \* ⟩ := |+0+⟩ + |−1−⟩ (Can be generated with a success rate of  $1/32$  linear-optically [7])
- $\blacktriangleright$  **Resource cost**  $Q$ : Average number of  $|G^{(3)}_{*}\rangle$ 's required to successfully generate
- one  $|G\rangle$  state through post-selection
- 1. Simplify the graph of the desired graph state via **unraveling**.
- 2. Construct a **fusion network**.
- 3. Determine the **fusion order** and calculate the **resource cost**  $Q$ .
- 4. Iterate 1—3 a sufficient number of times and select the best one.

#### **STAGE 2: CONSTRUCTION OF FUSION NETWORK**

‣ **Type-II fusion** [6]: Measuring {*X* ⊗ *Z*, *Z* ⊗ *X*}

#### **STAGE 3: DETER**

**EXAMPLE**

# **NUMERICAL RESULTS**





#### **REFERENCES**

[1] R. Raussendorf et al., Ann. Phys. **321**, 2242 (2006). [2] R. Raussendorf et al., New J. Phys. **9**, 199 (2007). [3] S. Bartolucci et al., Nat. Commun. 14, 912 (2023). [4] M. Zwerger et al., Phys. Rev. A 85, 062326 (2012). [5] N. Shettel and D. Markham, Phys. Rev. Lett. 124, 110502 (2020). [6] D. E. Browne and T. Rudolph, Phys. Rev. Lett. 95, 010501 (2005). [7] M. Vernava, Phys. Rev. Lett. 100, 060502 (2008). [8] E. L. Lawler, SIAM J. Comput. 9, 558-565 (1980).



 $\blacktriangleright$  Finding all non-overlapping bipartitely-complete subgraphs  $\rightarrow O\left(\mid V \!\mid\! d_{\max}^4\right)$  $\triangleright$  Finding all non-overlapping cliques [8]  $\rightarrow O(\text{poly}(N_{\text{clique}}))$ 













=

 $(1 - \eta)$ 

2

2

